## Validated Constraint Solving Practicalities, Pitfalls, and New Developments

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This talk will:

- Review three filtering schemes.
- Explain rigor in linear relaxations.
- Give a sequence of examples where

1. basic constraint propagation fails but not interval Newton narrowing;
2. interval Newton narrowing fails but linear relaxations do not.

- Describe recent implementation and numerical experiments with our GlobSol system.
on three narrowing strategies


## General Problem

minimize $\varphi(x)$
subject to:

$$
\begin{gathered}
c_{i}(x)=0, i=1, \ldots, m_{1}, \\
g_{i}(x) \leq 0, i=1, \ldots, m_{2},
\end{gathered}
$$

where $\varphi: \boldsymbol{x} \rightarrow \mathbb{R}$ and $c_{i}, g_{i}: \boldsymbol{x} \rightarrow \mathbb{R}$,
and where $\boldsymbol{x} \subset \mathbb{R}^{n}$ is the
hyperrectangle (box) defined by

$$
\underline{x}_{i_{j}} \leq x_{i_{j}} \leq \bar{x}_{i_{j}}, 1 \leq j \leq m_{3},
$$

$i_{j}$ between 1 and $n$, where the $\underline{x}_{i_{j}}$ and $\bar{x}_{i j}$ are constant bounds.
If $\varphi$ is constant or absent, this problem becomes a general constraint problem; if, in addition $m_{2}=m_{3}=0$, this problem becomes a nonlinear system of equations.

## Basic Constraint Propagation

Our view:

- We begin with bounds on the variables.
- We solve a constraint for a variable $x_{i}$, obtaining $x_{i}=g\left(x_{1}, \ldots, x_{i-1}, x_{i+1}, \ldots x_{n}\right)$.
- We use the bounds on $x_{j}, j \neq i$ to obtain (hopefully) narrower bounds on $x_{i}$ (say by evaluating $g$ with interval arithmetic).
- We "propagate" the new bounds on $x_{i}$, that is, we use the new bounds on $x_{i}$ in the constraints in which it occurs to obtain new bounds on the other variables.


## Basic Constraint Propagation

## An Example

Take the constraint system

$$
\begin{aligned}
& c_{1}(x)=x_{1}^{2}-2 x_{2}, \quad c_{2}(x)=x_{2}^{2}-2 x_{1}, \\
& x_{1} \in[-1,1], \quad x_{2} \in[-1,1] .
\end{aligned}
$$

1. Solve for $x_{2}$ in $c_{1}$, to obtain $x_{2}=x_{1}^{2} / 2$, then plug $x_{1}=[-1,1]$ into $x_{1}^{2} / 2$, to obtain $x_{2} \in[0,0.5]$.
2. Solve $c_{2}$ for $x_{1}$ to obtain $x_{1}=x_{2}^{2} / 2$, then plug the narrower value of $x_{2}$ into $x_{2}^{2} / 2$, to obtain $x_{1} \in[0,0.125]$.
3. Use $c_{1}$ again to obtain an even narrower value for $x_{2}$.
4. This process can be continued to convergence to $x_{1}=0, x_{2}=0$.

## Basic Constraint Propagation

When it does and does not work

- Basic constraint propagation only works for linear systems when a permutation of the rows and columns leads to a diagonally dominant system.
- For nonlinear systems, the Jacobi matrix should be permutable to a diagonally dominant system, or so preconditioned.
- Possible research direction: try symbolic preconditioning. (See me for references on how.)


## Basic Constraint Propagation

Example when it does not work
Take the constraint system

$$
\begin{aligned}
& c_{1}(x)=x_{1}^{3}+x_{1}-x_{2}, \quad c_{2}(x)=-2 x_{1}-x_{2}, \\
& x_{1} \in[-.5, .5], \quad x_{2} \in[-.25, .25] .
\end{aligned}
$$

- There is a unique solution $c_{1}=0, c_{2}=0$ at $x_{1}=0$, $x_{2}=0$.
- Solving $c_{1}$ for $x_{2}$ as in the previous example gives $x_{2}=\left(x_{1}^{3}+x_{1}\right)$.
- Solving $c_{1}=0$ for $x_{2}$ and using the exact range of $\left(x_{1}^{3}+x_{1}\right)$ for $x_{1} \in[-.5, .5]$ gives $x_{2} \in[-.625, .625]$, no improvement.
- Solving for $x_{2}$ in the second equation gives the range of $-2 x_{1}$ over $x_{1} \in[-.5, .5]$ is $x_{2} \in[-1,1]$, also not an improvement.
- The only remaining alternatives are to solve for $x_{1}$ in $c_{1}$ or $c_{2}$. Solving for $x_{1}$ in $c_{2}$ gives no improvement, but solving for $x_{1}$ in $c_{1}$ and plugging in $x_{2} \in[-.25, .25]$ gives $x_{1} \in[-.237, .237]$, an improvement.
- Additional applications of the process give no additional narrowing.
on three narrowing strategies


## Interval Newton Narrowing

In the above example,

$$
F(x)=\binom{c_{1}(x)}{c_{2}(x)}=\binom{x_{1}^{3}+x_{1}-x_{2}}{-2 x_{1}-x_{2}},
$$

and an element-wise interval extension of the Jacobi matrix of $F$ over the initial $\boldsymbol{x}$ is

$$
\boldsymbol{F}^{\prime}(\boldsymbol{x}) \in\left(\begin{array}{cc}
{[1,1.75]} & -1 \\
-2 & -1
\end{array}\right) .
$$

If the inverse of the midpoint matrix for $\boldsymbol{F}^{\prime}(\boldsymbol{x})$ is used as a preconditioner matrix, then, if $\check{x}=(0,0)^{T}$, the preconditioned system becomes

$$
\left(\begin{array}{cc}
{[0.8888,1.1112]} & {[0,0]} \\
{[-0.2223,0.2223]} & {[1,1]}
\end{array}\right)\binom{v_{1}}{v_{2}}=\binom{0}{0}
$$

and, using the interval Gauss-Seidel method, new bounds for $v$ are

$$
v \in\binom{0}{0}
$$

that is, we obtain the solution sharply.

## When Interval Newton Narrowing Fails

Take the constraint system
$c_{1}(x)=x_{1}^{2}-x_{2}=0, \quad g_{1}(x)=x_{2}-x_{1} \leq 0$, $x_{1} \in[0,1], \quad x_{2} \in[0,1]$.

- One easily checks that basic constraint propagation fails for this problem.
- To obtain lower and upper bounds on this solution set, we may solve the corresponding constrained optimization problems with objective functions min $x_{1}$, $\max x_{1}, \min x_{2}$ and $\max x_{2}$, subject to $-x_{1} \leq-.5, x_{1} \leq .5,-x_{2} \leq-.5, x_{2} \leq .5$.


# When Interval Newton Narrowing Fails 

## (continued)

- The Fritz-John equations for the problem with min can be written as

$$
\left(\begin{array}{c}
u_{0}-u_{1}-u_{2}+u_{3}+2 x_{1} v_{1} \\
u_{1}-u_{4}+u_{5}-v_{1} \\
u_{1}\left(x_{2}-x_{1}\right) \\
-u_{2} x_{1} \\
u_{3} x_{1} \\
-u_{4} x_{2} \\
u_{5} x_{2} \\
x_{1}^{2}-x_{2} \\
u_{0}+u_{1}+u_{2}+u_{3}+u_{4}+u_{5}+v_{1}^{2}-1
\end{array}\right)
$$

- Using $u_{i} \in[0,1], v_{1} \in[-1,1]$ for the Lagrange multipliers, the interval Jacobi matrix for this system contains many singular matrices, and the corresponding interval Newton method thus cannot succeed.


## Linear Relaxations

The basic idea

- If the objective $\varphi$ is replaced by linear function $\varphi^{(\ell)}$ such that $\varphi^{(\ell)}(x) \leq \varphi(x)$ for $x \in \boldsymbol{x}$, then the resulting problem has global optimum less than or equal to the global optimum of the original problem.
- If each inequality constraint $g_{i}$ replaced by a linear function $g_{i}^{(\ell)}$ such that $g_{i}^{(\ell)}(x) \leq g_{i}(x)$ for $x \in \boldsymbol{x}$, then the resulting problem, has optimum that is less than or equal to the optimum of the original problem.
- If there are equality constraints, then each equality constraint can be replaced by two linear inequality constraints, and these inequality constraints can be replaced as above by linear inequality constraints.
- The resulting linear program is termed a linear relaxation.


## Linear Relaxations

## Our Previous Example

$c_{1}(x)=x_{1}^{2}-x_{2}=0, \quad g_{1}(x)=x_{2}-x_{1} \leq 0$, $x_{1} \in[0,1], \quad x_{2} \in[0,1]$.

- Lower bounds of a convex function are tangent lines and upper bounds are secant lines.
- A corresponding linear program for computing an upper bound on $x_{2}$, using two underestimators for the convex function $x_{2}=x_{1}^{2}$, is:
minimize $-x_{2}$
subject to

$$
\begin{aligned}
& x_{2} \leq x_{1}(\text { the overestimator }) \\
& x_{2} \geq .125+.5\left(x_{1}-.25\right), \\
& x_{2} \leq x_{1}(\text { the original constraint }), \\
& x_{1} \in[0,1], x_{2} \in[0,1]
\end{aligned}
$$

## Linear Relaxation Example

## (continued)

- The exact minimum to this linear program is $\varphi=-.5$, corresponding to $x_{2} \leq 0.5$.
- Thus, we have narrowed $x_{2}$ to $x_{2} \in[0,0.5] \subset[0,1]$.
- Basic constraint propagation now converges.


## Rigor in Linear Relaxations

1. Typical procedures have been to compute the coefficients of the linear relaxation with floating point arithmetic, then to solve the relaxation with a state-of-the-art LP solver.

2 . With carefully considered directed rounding and interval arithmetic, we can form a machine-representable LP that is an actual relaxation of the original problem.
3. Neumaier and Shcherbina, as well as Jansson, have presented a simple technique to utilize the duality gap to obtain a rigorous lower bound on the solution to an LP, given approximate values of the dual variables.
4. Combining (2) and (3) gives a procedure for rigorous computations of lower bounds on the solution to the original problem.

## Implementation in GlobSol

- We have implemented linear relaxations in GlobSol.
- Initial experiments indicate the technique makes possible solution of problems that were previously intractable within GlobSol.
- A preprint of experimental results is available.
- GlobSol still is not fully competitive with other packages using relaxations in a non-validated way (e.g. BARON).
- One possibility for improvement: Use a better LP solver. (GlobSol presently is using a free one from the SLATEC library.)

