## On Existence and Uniqueness Verification for Non-Smooth Functions

by

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- We will show actual computations, to illustrate the relationship between traditional interval Newton methods and degree theory.
- We will illustrate how the computations can succeed or break down in non-smooth problems.

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#### The General Question

Let F(x) = 0 represent a system of nequations in n unknowns, and suppose  $\check{x}$  is a numerical approximation to a solution  $x^*$ ,  $F(x^*) = 0$ . We wish to compute bounds

$$oldsymbol{x} \;=\; (oldsymbol{x}_1, oldsymbol{x}_2, \dots, oldsymbol{x}_n) \ =\; ([\underline{x}_1, \overline{x}_1], [\underline{x}_2, \overline{x}_2], \dots, [\underline{x}_n, \overline{x}_n],$$

such that  $\check{x}$  is the center of  $\boldsymbol{x}$ , and such that  $\boldsymbol{x}$  is guaranteed to contain a solution  $x^*$  to F(x) = 0. That is,

Given 
$$F : \boldsymbol{x} \to \mathbb{R}^n$$
, where  $\boldsymbol{x} \in \mathbb{IR}^n$ ,  
*rigorously* verify:  
• there exists a  $x^* \in \boldsymbol{x}$  such that  
 $F(x^*) = 0.$ 
(1)

Here,  $\mathbb{IR}^n$  represents the set of interval n-vectors.

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#### **Interval Newton Methods**

The Traditional Setting

If the Jacobi matrix  $F'(x^*)$  is non-singular and continuous in  $\boldsymbol{x}$ , then we can use an interval Newton method:

$$\tilde{\boldsymbol{x}} = \boldsymbol{N}(F; \boldsymbol{x}, \check{x}) = \check{x} + \boldsymbol{v},$$

where

$$\Sigma(\boldsymbol{A}, -F(\check{x})) \subset \boldsymbol{v},$$

where  $\boldsymbol{A}$  is a Lipschitz matrix for F over  $\boldsymbol{x}$ ,

and where 
$$\Sigma(\mathbf{A}, -F(\check{x}))$$
  
= { $x \in \mathbb{R}^n \mid \exists A \in \mathbf{A}$  with  $AX = -F(\check{x})$  }.

We have:

**Theorem 1** (see Neumaier's book) Suppose  $\tilde{\boldsymbol{x}} = \boldsymbol{N}(F; \boldsymbol{x}, \check{\boldsymbol{x}})$  is the image of  $\boldsymbol{x}$  and  $\check{\boldsymbol{x}}$ under an interval Newton method. If  $\tilde{\boldsymbol{x}} \subseteq \boldsymbol{x}$ , it follows that there exists a unique solution of  $F(\boldsymbol{x}) = 0$  within  $\boldsymbol{x}$ .

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### Modifications for Singular / Non-Smooth Systems

The Topological Degree

- We can verify existence of solutions to F(x) = 0 within  $\boldsymbol{x}$ , even when  $\det(F'(x^*)) = 0$ .
- We do this with the *topological degree*  $d(F, \boldsymbol{x}, 0)$  of F over  $\boldsymbol{x}$ .
- If  $\det(F'(x)) \neq 0$  when F(x) = 0, then  $d(F, \boldsymbol{x}, 0) = \sum_{\substack{x \in \operatorname{int}(\boldsymbol{x}), \\ F(x) = 0}} \operatorname{sgn}(\det(F'(x))).$
- The integer d(F, x, 0) is continuous in F and depends only on values of F on the boundary ∂x, so F' may be singular or non-smooth in the interior int(x).
- $d(F, \boldsymbol{x}, 0) \neq 0 \Rightarrow F(x) = 0$  has a solution in  $x^* \in \boldsymbol{x}$ .

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#### Modifications for Singular / Non-Smooth Systems

The Theorem Used in the Algorithms

• The boundary of  $\boldsymbol{x}$  consists of:

$$egin{aligned} oldsymbol{x}_{\overline{k}} &\equiv \left(oldsymbol{x}_{1},\ldots,oldsymbol{x}_{k-1},\overline{x}_{k},oldsymbol{x}_{k+1},\ldots,oldsymbol{x}_{n}
ight)^{T},\ oldsymbol{x}_{\overline{k}} &\equiv \left(oldsymbol{x}_{1},\ldots,oldsymbol{x}_{k-1},\overline{x}_{k},oldsymbol{x}_{k+1},\ldots,oldsymbol{x}_{n}
ight)^{T},\ \mathrm{where} \;k=1,\ldots,n. \end{aligned}$$

- For fixed  $\ell$ ,  $1 \leq \ell \leq n$ , define  $F_{\neg \ell}(x) = (f_1(x), \dots, f_{\ell-1}(x), f_{\ell+1}(x), \dots, f_n(x))^T$ .
- For this  $\ell$ , define  $\underline{K_0(s)}$  as that subset of  $\{k | k \in \{1, \ldots, n\}\}$  such that  $F_{\neg \ell} = 0$  has solutions on  $\boldsymbol{x}_{\underline{k}}$  and  $\operatorname{sgn}(f_{\ell}) = s$  at these solutions; similarly define  $\overline{K_0(s)}$  such that  $F_{\neg \ell} = 0$  has solutions on  $\boldsymbol{x}_{\overline{k}}$  and  $\operatorname{sgn}(f_{\ell}) = s$  at these solutions, where  $s \in \{-1, +1\}.$

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#### The Theorem Used in the Algorithms (continued)

**Theorem 2** If F is continuous,  $F \neq 0$  on  $\partial \boldsymbol{x}$ , and there is an  $\ell$ ,  $1 \leq \ell \leq n$ , such that: (1)  $F_{\neg \ell} \neq 0$  on  $\partial \boldsymbol{x}_{\underline{k}}$  or  $\partial \boldsymbol{x}_{\overline{k}}$ ,  $k = 1, \ldots, n$ ; (2)  $det(F'_{\neg \ell}) \neq 0$  whenever  $F_{\neg \ell} = 0$  on  $\partial \boldsymbol{x}$ . Then

$$d(F, \boldsymbol{x}, 0) = (-1)^{\ell-1} s$$

$$\cdot \begin{cases} \sum_{k \in \underline{K}_0(s)} (-1)^k \\ \sum_{x \in \boldsymbol{x}_k} \operatorname{sgn} & \left| \frac{\partial F_{\neg \ell}}{\partial x_1 x_2 \dots x_{k-1} x_{k+1} \dots x_n} (x) \right| \\ + \sum_{k \in \overline{K}_\tau(s)} (-1)^{k+1} \end{cases}$$

$$\sum_{\substack{x \in \mathbf{X}_{\overline{k}} \\ F_{\neg \ell}(x) = 0}} \operatorname{sgn} \left| \frac{\partial F_{\neg \ell}}{\partial x_1 x_2 \dots x_{k-1} x_{k+1} \dots x_n} (x) \right| \right\}.$$

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## Simplifications to Make It Practical

In our methods, we

- 1. precondition F;
- 2. choose the coordinate widths  $w(\boldsymbol{x}_k)$ ,  $1 \leq k \leq n$  to have  $\boldsymbol{F}_{\neg \ell}(\boldsymbol{x}_k) \neq 0$  and  $\boldsymbol{F}_{\neg \ell}(\boldsymbol{x}_k) \neq 0$  for all k except k = n - p to k = n, where p is the dimension of the null space. This eliminates most terms in Theorem 2.
- 3. We then use a *p*-dimensional search on the remaining several faces of  $\boldsymbol{x}$  to find the solutions of  $F_{\neg \ell} = 0$ .
- 4. In certain instances, we use a heuristic to guess the value of  $d(F, \boldsymbol{x}, 0)$ .

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### Simplifications to Make It Practical

The Preconditioner

We use incomplete LU factorization (with full pivoting) to put the Jacobi matrix into the form

$$YF'(x^*) \approx \begin{pmatrix} 1 & 0 & \dots & 0 & * \\ 0 & 1 & 0 & \dots & 0 & * \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & \dots & 0 & 1 & * \\ 0 & \dots & 0 & 0 & 0 \end{pmatrix}$$

(for the case where the rank defect is 1).

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#### An Example

$$f_1(x) = x_1 + x_2 + x_3, f_2(x) = -x_2 + x_3^3, f_3(x) = x_2 + x_3^3,$$

with

$$\boldsymbol{x} = ([-0.02, 0.02], [-0.01, 0.01], [-0.01, 0.01])^T.$$
$$\boldsymbol{F}'(\boldsymbol{x}) = \begin{pmatrix} 1 & 1 & 1 \\ 0 & -1 & [0, .0003] \\ 0 & 1 & [0, .0003] \end{pmatrix}, \boldsymbol{Y} = \begin{pmatrix} 1 & 0 & -1 \\ 0 & -1 & 0 \\ 0 & 1 & 1 \end{pmatrix}.$$

We thus have

$$YF(x) = \begin{pmatrix} 1 & 0 & [0.9997, 1] \\ 0 & 1 & [-0.0003, 0] \\ 0 & 0 & [0, 0.0006] \end{pmatrix},$$
$$YF(x) \approx \begin{pmatrix} x_1 + x_3 - x_3^3 \\ x_2 - x_3^3 \\ 2x_3^3 \end{pmatrix}.$$

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- We will find solutions to  $F_{\neg 3} = 0$  on the boundary of  $\boldsymbol{x}$  at which  $\operatorname{sgn}(f_3) = +1$ .
- We choose the widths of  $\boldsymbol{x}$  appropriately, then we use mean value extensions to show

$$-f_1 \neq 0$$
 on  $\boldsymbol{x}_1$  and  $\boldsymbol{x}_{\overline{1}}$  and

- $-f_2 \neq 0$  on  $\boldsymbol{x}_{\underline{2}}$  and  $\boldsymbol{x}_{\overline{2}}$ .
- We then proceed with the interval Gauss–Seidel method on  $x_{\underline{3}}$  and  $x_{\overline{3}}$ .

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• No solutions on  $x_1$ :

$$(YF)_1(\boldsymbol{x}_{\underline{1}}) \subseteq (YF)_1(0,0,0) + 1 \cdot (-0.02) \\ + [0.9997,1] \cdot [-0.01,0.01] \\ \subseteq [-0.03,-0.01].$$

- Similarly, on  $\boldsymbol{x}_{\overline{1}}$ :  $(Yf)_1(\boldsymbol{x}_{\overline{1}}) \subseteq [.01, .03].$
- We thus have verified  $(YF)_{\neg 3} \neq 0$  on  $\boldsymbol{x}_{\underline{1}} = (-0.02, [-0.01, 0.01], [-0.01, 0.01])^T$ and  $\boldsymbol{x}_{\overline{1}} = (+0.002, [-0.01, 0.01], [-0.01, 0.01])^T$ .
- Similarly, we use mean value extensions for  $(Yf)_2$  on  $\boldsymbol{x}_2$  and  $\boldsymbol{x}_{\overline{2}}$  to verify that  $(YF)_{\neg 3} \neq 0$  on  $\boldsymbol{x}_2$  and  $\boldsymbol{x}_{\overline{2}}$ .

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Verifying solutions on  $x_{\underline{3}}$  and  $x_{\overline{3}}$ 

For  $x_{\underline{3}}$ :

• Plug in  $x_3 = -0.01$  and apply the interval Gauss–Seidel method. Setting

{Mean Value Extension for  $(YF)_1$ } = 0 gives

$$(YF)_1(0, 0, -.01) + 1 \cdot (x_1 - 0) + 0 \cdot x_2 + [0.9997, 1](0) = 0;$$

solving this for  $x_1$  gives

$$x_1 \in 0 - \{(YF)_1(0, 0, -0.01) \\ -([-0.01, -0.009997] \cdot 0)\} / 1 \\ = 0.009999.$$

• Similarly,  $x_2 \in [-10^{-6}, -10^{-6}].$ 

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Verifying solutions on  $x_{\underline{3}}$  and  $x_{\overline{3}}$ 

• Thus,

 $x \in ([0.009999], [-10^{-6}], [-0.01])^T = x^{(1)}.$ 

- An interval evaluation of  $F(\boldsymbol{x}^{(1)})$  gives  $YF(x) \in ([0, 0], [0, 0], [-2 \times 10^{-6}, -2 \times 10^{-6}])^T$
- Since  $(YF)_3 < 0$ , this solution can be ignored.
- Similar computations on  $\boldsymbol{x}_{\overline{3}}$  give a single point  $\boldsymbol{x}^{(\overline{1})}$  at which  $(YF)_3(\boldsymbol{x}^{(\overline{1})}) > 0$  and at which which

$$\det\left(\frac{\partial (YF)_{\neg 3}}{\partial x_1 \partial x_2}(\boldsymbol{x}^{\overline{1}})\right) > 0$$

• Combining these facts into the sum in the theorem gives a topological degree  $d(YF, \boldsymbol{x}, 0) = 1.$ 

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Summary

- $d(YF, \boldsymbol{x}, 0) = 1$  proves existence of a solution in  $\boldsymbol{x}$ .
- This result has been proven with 2 \* (n 1) mean-value-extension evaluations of (n 1) components of F and with two incomplete Gauss-Seidel sweeps.
- Can the same process succeed when the components of *F* are non-smooth?

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# The Principles Behind Our Degree Computation

- For the interval Gauss–Seidel method, we "solve" the preconditioned variable for the *i*-th variable in the *i*-th equation.
- Success depends on the magnitude of the off-diagonal elements being larger than the mignitude of the diagonal elements of  $Y \mathbf{F}'(\mathbf{x})$ .
- In degree computation, we fix the last variables, eliminating the uncontrolled widths in the last column of  $Y \mathbf{F}'(\mathbf{x})$ .

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#### A Non-Smooth Example

Define

$$f_1(x) = x_1 + x_2 + x_3,$$
  

$$f_2(x) = \begin{cases} -x_2 + x_3^3 & \text{if } x_2 \ge 0, \\ -5x_2 + x_3^3 & \text{if } x_2 < 0, \end{cases}$$
  

$$f_3(x) = \begin{cases} x_2 + x_3^3 & \text{if } x_2 \ge 0, \\ 0.1x_2 + x_3^3 & \text{if } x_2 < 0, \end{cases}$$

and take

 $\boldsymbol{x} =$ 

 $([-0.02, 0.02], [-0.01, 0.01], [-0.01, 0.01])^T.$ Then

$$Y \mathbf{F}'(\mathbf{x}) = \begin{pmatrix} 1 & [-0.66\overline{6}, 0.66\overline{6}] & [1.0000, 1.0001] \\ 0 & [0.33\overline{3}, 1.66\overline{6}] & [-0.0001, 0] \\ 0 & [-0.81\overline{6}, 0.81\overline{6}] & [0, 0.000355] \end{pmatrix}.$$

• In this case, the off-diagonal entries of  $Y \mathbf{F}'$ (excluding the last column) are sufficiently small and narrow to allow the verification process to succeed:  $d(YF, \mathbf{x}, 0) = 1$ .

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# A Second Non-Smooth Example

$$f_1(x) = \begin{cases} x_1 + x_2 + x_3 & \text{if } x_2 \ge 0, \\ x_1 + 10x^2 + x_3 & \text{if } x_2 < 0, \end{cases}$$
  
$$f_2(x) = \text{same as in the previous example,}$$
  
$$f_3(x) = \text{same as in the previous example.}$$

In this case,

$$Y \mathbf{F}'(\mathbf{x}) = \begin{pmatrix} 1 & [-8.16\overline{6}, 8.16\overline{6}] & [1.0000, 1.00055] \\ 0 & [0.33\overline{3}, 1.66\overline{6}] & [-0.0001, 0] \\ 0 & [-0.81\overline{6}, 0.81\overline{6}] & [0, 0.000355] \end{pmatrix},$$

and the off-diagonal entries (excluding the last column) are <u>not</u> sufficiently small and narrow to allow the verification process to succeed.

• Details of these two examples can be found in

#### http://interval.louisiana.edu/ preprints/nonsmooth\_degree.pdf

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# **Additional Thoughts**

- In these algorithms, the effect of the lack of smoothness is similar to the effect of non-smoothness on traditional interval Gauss–Seidel methods for verification of non-singular zeros.
- The degree  $d(F, \boldsymbol{x}, 0)$  depends only on the values F on the boundary of  $\boldsymbol{x}$ .
  - The formula in the theorem also only involves values on the boundary.
  - In principle, there is no problem applying the theorem for successful verification, as long as F is smooth on the boundary.

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# Additional Thoughts (continued)

- In practice, non-smoothness inside  $\boldsymbol{x}$ makes simplifications as we have illustrated impossible in general.
- Direct application of the theorem involves (2n) global optimization problems.
- For details, see http://interval.louisiana.edu/ preprints/nonsmooth\_degree.pdf
- Our fast method may work anyway (as illustrated with our first non-smooth example).
- There may be other simplifications, for particular non-smooth cases, that we haven't yet discovered.

These transparencies will be available from http://interval.louisiana.edu/
 preprints.html

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